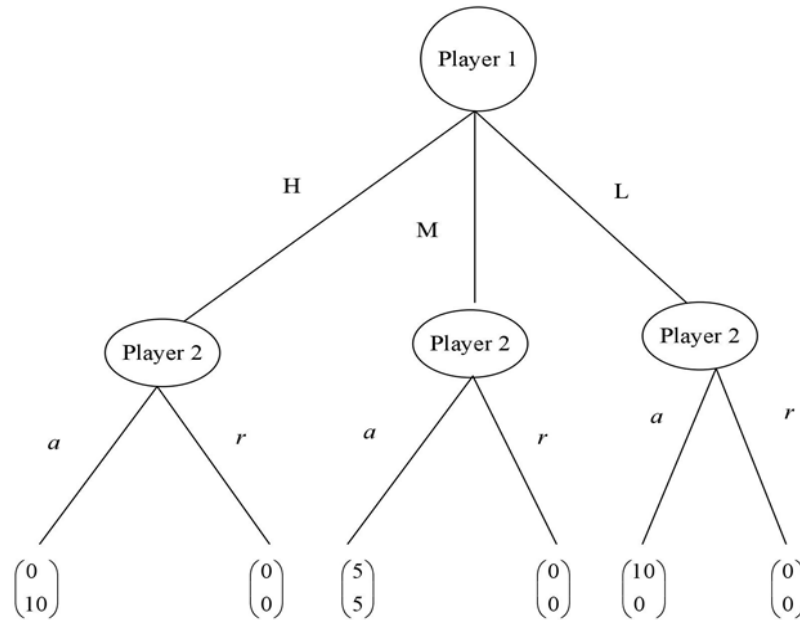


Homework # 5 [Due on Tuesday April 14th, 2026]

1. Find a story that can be represented by an extensive form game. Identify the: (i) set of players, (ii) set of actions, (iii) time structure of the game and (iv) payoffs.

- For instance, the example of "Antz vs A Bug's life"

2. Consider the following extensive form game.



(a) What are the strategies for player 1?

- The strategies for player 1 are H , M , and L .

(b) What are the strategies for player 2?

- The strategies for player 2 are: $\{aaa, aar, arr, rrr, rra, raa, ara, rar\}$

(c) Take your results from a) and b) and construct a matrix representing its normal form game representation.

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		Player 2							
		aaa	aar	arr	rrr	rra	raa	ara	rar
Player 1	H	0, 10	0, 10	0, 10	0,0	0,0	0,0	0,10	0,0
	M	5, 5	5, 5	5, 5	0,0	0,0	5,5	0,0	5,5
	L	10, 0	10, 0	10, 0	0,0	10,0	10,0	10,0	0,0

3. Consider the following normal form game

		Player 2	
		<i>L</i>	<i>R</i>
Player 1	<i>U</i>	-10, 10	0, 12
	<i>C</i>	-12, 0	2, 2
	<i>D</i>	-9, 0	1, 0

- (a) Find strictly dominant strategies (if any) for player 1 and for player 2.
- There are no strictly dominant strategies for both players 1 and 2.
- (b) Find strictly dominated strategies (if any) for player 1 and for player 2.
- **For player 1**, strategy *U* is strictly dominated because when player 2 chooses strategy *L*, he could maximize his payoff by choosing strategy *D*, and similarly, by choosing strategy *C* when player 2 chooses strategy *R*. **For player 2**, neither strategy *L* nor *R* is strictly dominated because when player 1 chooses strategy *D*, player 2 obtains the same payoff by either choosing strategy *L* or *R*, both of which give him exactly the same payoff of 0.
- (c) If you apply iterative deletion of strictly dominated strategies (IDSDS), what is the surviving strategy pair (or pairs)? Explain the steps you use in IDSDS, and why you use them.
- Since strategy *U* is strictly dominated, we can remove this strategy for player 1, resulting in a reduced-form matrix as follows:

		Player 2	
		<i>L</i>	<i>R</i>
Player 1	<i>C</i>	-12, 0	2, 2
	<i>D</i>	-9, 0	1, 0

At this stage of the game, none of the strategies are strictly dominated by one another, and therefore, the strategy profile that survives IDSDS becomes

$$\{(C, L), (C, R), (D, L), (D, R)\}$$

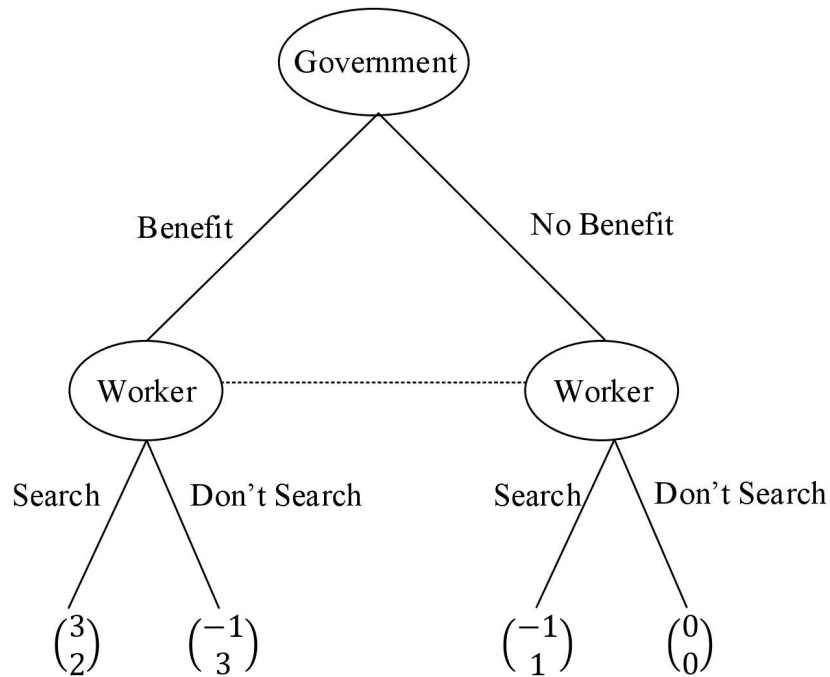
4. Consider the following simultaneous-move game between the government (row player), which decides whether to offer unemployment benefits, and an unemployed worker (column player), who chooses whether to search for a job. As you interpret from the payoff matrix below, the unemployed worker only finds it optimal to search for a job

when he receives no unemployment benefit; while the government only finds it optimal to help the worker when he searches for a job.

		Worker	
		<i>Search</i>	<i>Don't Search</i>
Government	<i>Benefit</i>	3, 2	-1, 3
	<i>No benefit</i>	-1, 1	0, 0

- (a) Represent this game in its extensive form (game tree), where the government acts first and the worker responds without observing whether the government offered unemployment benefits.

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- (b) Does government has strictly dominant strategies? How about the worker?
- There are no strictly dominant strategies for both government and worker. Specifically, when the worker chooses (not) to search for a job, the government will be better off (not) offering unemployment benefits. Whereas, when the government chooses (not) to offer unemployment benefits, the worker will be better off (not) searching for a job.
- (c) Find which strategy profile (or profiles) survive the application of IDSDS.
- In this context, every strategy survives IDSDS, which contribute to a strategy profile of

$\{(Benefit, Search), (Benefit, Don'tSearch), (NoBenefit, Search), (NoBenefit, Don'tSearch)\}$

5. Consider the following simultaneous-move game played by player 1 (in rows) and player 2 (in columns).

		<i>Player 2</i>		
		<i>x</i>	<i>y</i>	<i>z</i>
<i>Player 1</i>	<i>a</i>	2, 3	1, 4	3, 2
	<i>b</i>	5, 1	2, 3	1, 2
	<i>c</i>	3, 7	4, 6	5, 4
	<i>d</i>	4, 2	1, 3	6, 1

- (a) Which strategy pairs survive the application of iterative deletion of strictly dominated strategies (IDSDS)?
- For player 2 (column player), strategy z is strictly dominated by y . We can then delete column z , leaving us with the following reduced-form matrix.

		<i>Player 2</i>	
		<i>x</i>	<i>y</i>
<i>Player 1</i>	<i>a</i>	2, 3	1, 4
	<i>b</i>	5, 1	2, 3
	<i>c</i>	3, 7	4, 6
	<i>d</i>	4, 2	1, 3

For player 1 (row player), strategy a is strictly dominated by b . After deleting the row corresponding to a , we obtain

		<i>Player 2</i>	
		<i>x</i>	<i>y</i>
<i>Player 1</i>	<i>b</i>	5, 1	2, 3
	<i>c</i>	3, 7	4, 6
	<i>d</i>	4, 2	1, 3

At this point, we cannot delete any more strategies for players 1 or 2 if we restrict them to use pure strategies. However, if we allow player 1 to randomize between the strategies that provide the highest payoff, b and c . In particular, assigning a probability p to strategy b and the remaining probability $1 - p$ to strategy c , player 1's expected payoff when player 2 chooses strategy x (in

the left-hand column of the above matrix) is

$$5p + 3(1 - p) = 2p + 3$$

which is larger than player 1's payoff from strategy d , 4, as long as $2p + 3 > 4$, or solving for p , if $p > \frac{1}{2}$. Similarly, when player 2 chooses strategy y (in the right-hand column of the above matrix), player 1's expected payoff from randomizing between b and c becomes

$$2p + 4(1 - p) = 4 - 2p$$

which is larger than player 1's payoff from strategy d , 1, as long as $4 - 2p > 1$, or solving for p , if $p < \frac{3}{2}$. This condition holds by assumption since probability p must be a number between 0 and 1. Therefore, any randomization between strategies b and c that assigns more than 50% probability on strategy b (that is, $p > 1/2$) yields a expected utility larger than the utility player 1 receives from strategy d . We can therefore claim that strategy d is strictly dominated, and delete the bottom row of the above matrix, leaving us with the followed reduced-form matrix.

		<i>Player 2</i>	
		<i>x</i>	<i>y</i>
<i>Player 1</i>	<i>b</i>	5, 1	2, 3
	<i>c</i>	3, 7	4, 6

At this point, we cannot delete any further strategies for players 1 or 2. Then, the strategy profiles surviving IDSDS are those in the four cells of the above matrix:

$$IDSDS = \{(b, x), (b, y), (c, x), (c, y)\}.$$

- (b) Using your results from part (a), show that there is no pure strategy Nash equilibrium (psNE) in this game.
- Using the strategy profiles that survived IDSDS, we can next underline best response payoffs, as depicted in the matrix below.

		<i>Player 2</i>	
		<i>x</i>	<i>y</i>
<i>Player 1</i>	<i>b</i>	<u>5</u> , 1	2, <u>3</u>
	<i>c</i>	3, <u>7</u>	<u>4</u> , 6

Since there is no cell where both players' payoffs are underlined, we can claim that there is no pure strategy Nash equilibrium in this game. There is, however, a mixed strategy Nash equilibrium, as we show in the next part of the exercise!

(c) Using your results from part (a), find a mixed strategy Nash equilibrium (msNE) in this game.

- *Player 1.* If player 1 is randomizing, he must be indifferent between pure strategies b and c . His expected utility from choosing b (in the top row of the above matrix) is

$$EU_1(b) = 5q + 2(1 - q) = 3q + 2$$

while his expected utility from selecting c (in the bottom row of the matrix) is

$$EU_1(c) = 5q + 4(1 - q) = 4 - q.$$

Then, player 1 is indifferent between b and c if and only if $EU_1(b) = EU_1(c)$, which implies that

$$3q + 2 = 4 - q$$

and, after rearranging, $4q = 2$, or $q = \frac{1}{2}$.

- *Player 2.* If player 2 is randomizing, he must be indifferent between his pure strategies x and y . His expected utility from choosing x (in the left-hand column of the above matrix) is

$$EU_2(x) = 1p + 7(1 - p) = 7 - 6p$$

while his expected utility from selecting y (in the right-hand column of the matrix) is

$$EU_2(y) = 3p + 6(1 - p) = 6 - 3p.$$

Then, player 2 is indifferent between x and y if and only if $EU_2(x) = EU_2(y)$, which implies that

$$7 - 6p = 6 - 3p$$

or, after rearranging, $1 = 3p$, or $p = \frac{1}{3}$.

- Therefore, the mixed strategy Nash equilibrium of the game is

$$\left\{ \left(\frac{1}{3}b, \frac{2}{3}c \right), \left(\frac{1}{2}x, \frac{1}{2}y \right) \right\}$$

where the first pair indicates player 1's randomization between b and c with

probabilities $1/3$ and $2/3$ respectively, while the second pair represents player 2's randomization between x and y , each with 50% probability.